

Teaching Commutative Algebra and Algebraic Geometry using Computer Algebra Systems

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Department of Mathematics,
Simon Fraser University
British Columbia, CANADA

Computer Algebra in Education
ACA 2012, Sofia, Bulgaria
June 25-28, 2012

MATH 441 Commutative Algebra and Algebraic Geometry Simon Fraser University, 2006, 2008, 2010, 02012

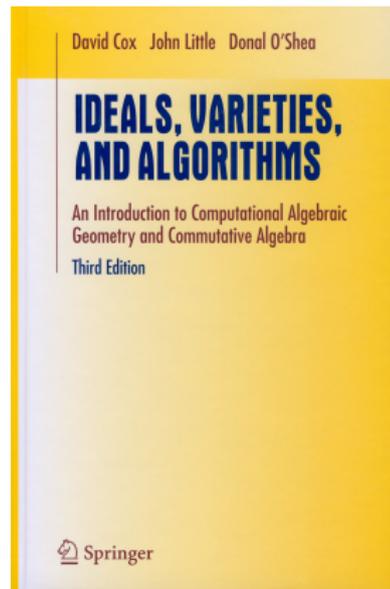
- Who takes the course?
- Textbook and course content.
- Maple and Assessment.
- **Three applications.**
 - Read material (paper).
 - Reproduce computational results.
 - Correct errors.
- Course project for graduate students.

Who takes the course?

4th undergraduate students and
1st year graduate students.

major	2006	2008	2010	2012	total
mathematics	5	15	11	27	58
computing	0	0	0	1	1
math & cmpt	0	2	2	2	6
other	0	4	0	0	4
graduate	5	6	4	1	16
total	10	27	17	31	85

- Ch. 1 Geometry, Algebra and Algorithms
 - Ch. 2 **Groebner Bases**
 - Ch. 3 Elimination Theory
 - Ch. 4 The Algebra-Geometry Dictionary
 - Ch. 5 Quotient Rings
 - Ch. 6 Automatic Geometric Theorem Proving
 - Ch. 7 Invariant Theory of Finite Groups
 - Ch. 8 Projective Algebraic Geometry
- Appendix C. Computer Algebra Systems



- 1 Varieties, [Graphing varieties](#), Ideals.
- 2 Monomial orderings and the division algorithm.
The Hilbert basis theorem.
Gröbner bases and [Buchberger's algorithm](#).
- 3 [Solving equations](#) (using Gröbner bases).
Elimination theory and resultants.
- 4 Hilbert's Nullstellensatz.
Radical ideals and [radical membership](#).
Zariski topology.
Irreducible varieties, [prime ideals](#), maximal ideals.
[Ideal decomposition](#).
- 5 Quotient rings, [computing in quotient rings](#).
- 6 [Applications](#) (of Gröbner bases).

Maple and Assessment

- One intro Maple tutorial in lab.
- Detailed examples worksheet for self study.
- Five in class demos.
- Maple worksheet handouts.

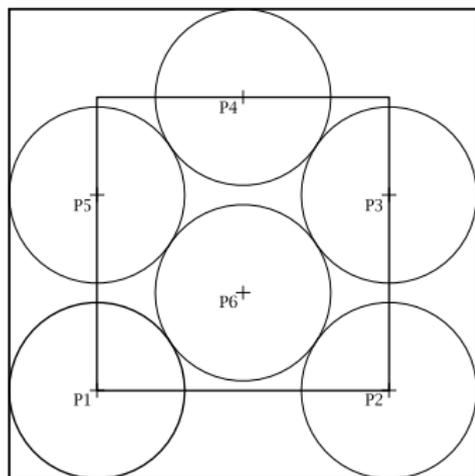
Maple and Assessment

- One intro Maple tutorial in lab.
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	MATH 441	MATH 819
6 assignments	60%	60%
project	–	10%
24 hour final	40%	30%

- Assignment questions, final exam, and project need Maple.
- Post take home final on web at 9am.
Hand in following day before 10am in person.

Application 1: Circle Packing Problems



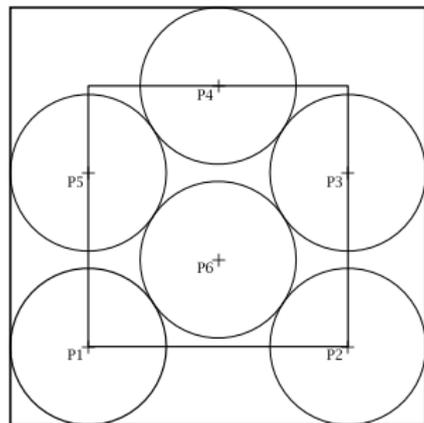
Pack $n = 6$ circles in the unit square maximizing the radius r .

Pack $n = 6$ points in the unit square maximizing their separating distance m .

$$r = \frac{m}{2(m+1)}$$

D. Würtz, M. Monagan and R. Peikert.
The History of Packing Circles in a Square.
MapleTech, Birkhauser, 1994.

Application 1: Circle Packing Problems



Given a packing, find m .

Let $P_i = (x_i, y_i)$ for $1 \leq i \leq 6$.

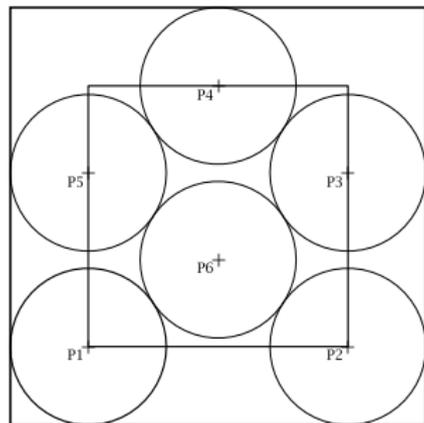
So $P_1 = (0, 0)$, $P_6 = (x_6, y_6)$, etc.

Pythagoras: $(x_6 - x_1)^2 + (y_6 - y_1)^2 = m^2$.

Symmetry: $x_6 = 1/2$, $y_6 = (y_0 + y_5)/2$.

Do not solve for $m, x_1, \dots, x_6, y_1, \dots, y_6$. Instead let

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Do not solve for $m, x_1, \dots, x_6, y_1, \dots, y_6$. Instead let

$I = \langle x_6 - \frac{1}{2}, x_6^2 + y_6^2 - m^2, \dots \rangle \subset \mathbb{Q}[x_1, \dots, x_6, y_1, \dots, y_6, m]$

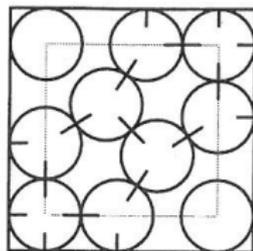
and compute a Gröbner basis G for $I \cap \mathbb{Q}[m] = \langle g \rangle$.

I get $G = \{(4m^2 - 5)(36m^2 - 13)\}$.

Figure out that $4m^2 - 5 = 0$ is a degenerate case.

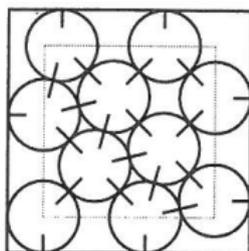
Application 1: Circle Packing Problems

Case $n = 10$



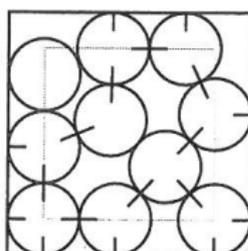
$m = 0.41953$

J. Schaer



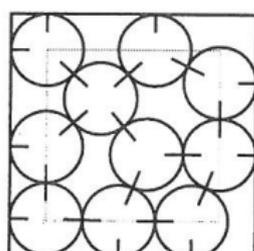
$m = 0.42013$

R. Milano



$m = 0.42118$

G. Valette



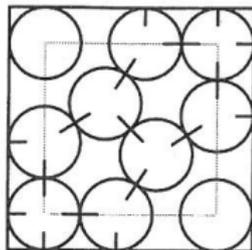
$m = 0.42129$

WMP

What can go wrong?

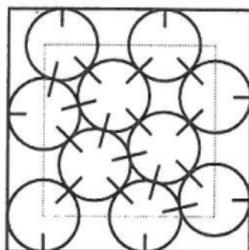
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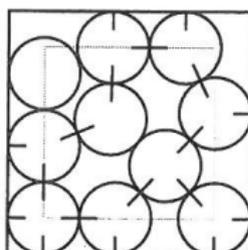
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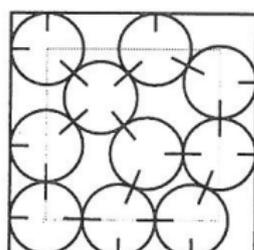
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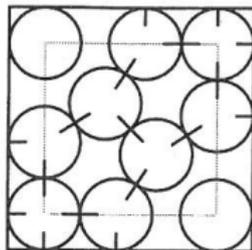
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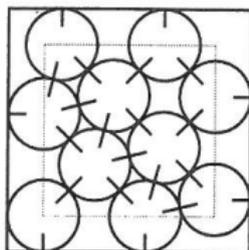
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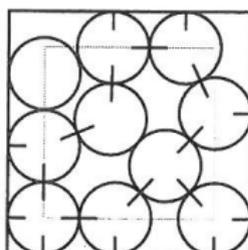
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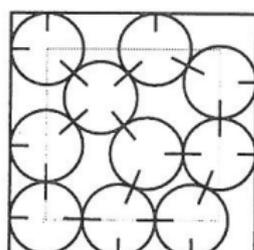
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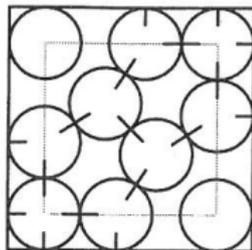
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Errors in the figures.

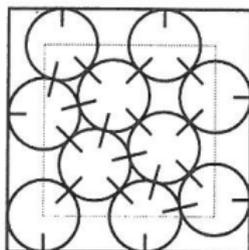
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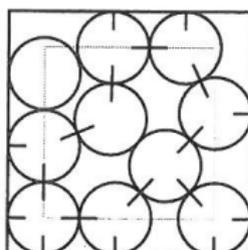
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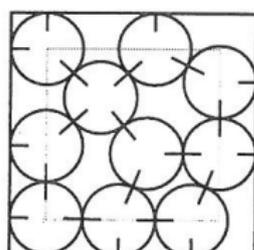
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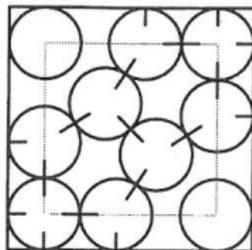
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Setup may have degenerate solutions.

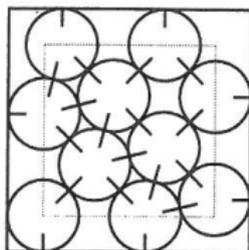
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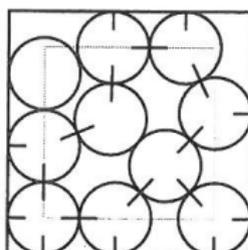
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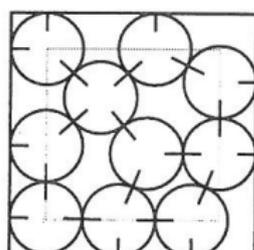
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What can go wrong?

Input equations incorrectly.

Errors in the figures.

Setup may have degenerate solutions.

Too many quadratic equations \implies long time.

Application 2: Graph Coloring and Hilbert's Nullstellensatz.

Which of these graphs can be colored with three colors?

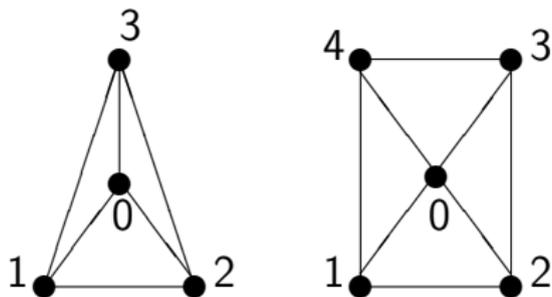
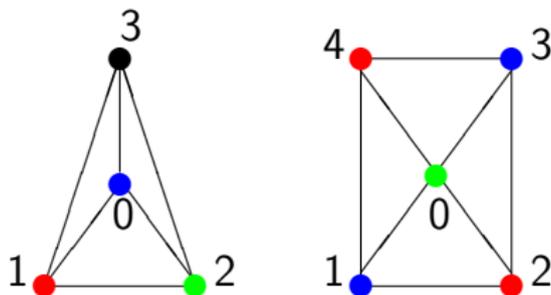


Figure: Wheel graphs W_3 and W_4 .



Application 2: Graph Coloring and Hilbert's Nullstellensatz.

To color a graph on n vertices with $k = 3$ colors, set

$$S := \{x_1^k = 1, x_2^k = 1, \dots, x_n^k = 1\}.$$

For each edge $(u, v) \in G$ set $S := S \cup \left\{ \frac{x_u^k - x_v^k}{x_u - x_v} = 0 \right\}$.

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Theorem

G is k -colorable $\iff S$ has solutions over \mathbb{C} .

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Theorem

G is k -colorable $\iff S$ has solutions over \mathbb{C} .

Nice! So let $I = \langle x_1^k - 1, \dots, \rangle$.

Compute a reduced Gröbner basis B for I .

Et voila! $B = \{1\} \iff G$ is not k -colorable.

But

Theorem

Graph k -colorability is NP-complete for $k \geq 3$.

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Graph k -colorability is NP-complete for $k \geq 3$.

J.A. de Loera, J. Lee, P.N. Malkin and S. Margulies.

Hilbert's Nullstellensatz and an algorithm for proving combinatorial infeasibility.

In *Proc. ISSAC 2008*, ACM Press, 197–206, 2008.

Application 2: Graph Coloring and Hilbert's Nullstellensatz.

Let $V = \mathbb{V}(x_1^k - 1, \dots, \dots)$ and $I = \langle x_1^k - 1, \dots, \rangle$.

Theorem

G is NOT k -colorable $\iff V = \emptyset \stackrel{HNS}{\iff} 1 \in I$.

Application 2: Graph Coloring and Hilbert's Nullstellensatz.

Let $V = \mathbb{V}(x_1^k - 1, \dots, \dots)$ and $I = \langle x_1^k - 1, \dots, \rangle$.

Theorem

G is NOT k -colorable $\iff V = \emptyset \stackrel{\text{HNS}}{\iff} 1 \in I$.

But if $I = \langle f_1, f_2, \dots, f_m \rangle \subset \mathbb{Q}[x_1, x_2, \dots, x_n]$ then

$$1 \in I \implies 1 = h_1 f_1 + h_2 f_2 + \dots + h_m f_m$$

for some h_1, h_2, \dots, h_m in $\mathbb{Q}[x_1, x_2, \dots, x_n]$.

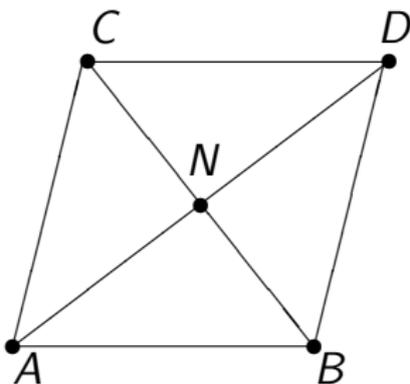
Idea 1: Try to find h_i with degree $d = 1, 2, 3, \dots$

Idea 2: The larger d the harder the combinatorial problem.

Idea 3: Replace \mathbb{Q} with \mathbb{F}_2 .

Get the students to experiment.

Application 3: Automatic Geometric Theorem Proving.

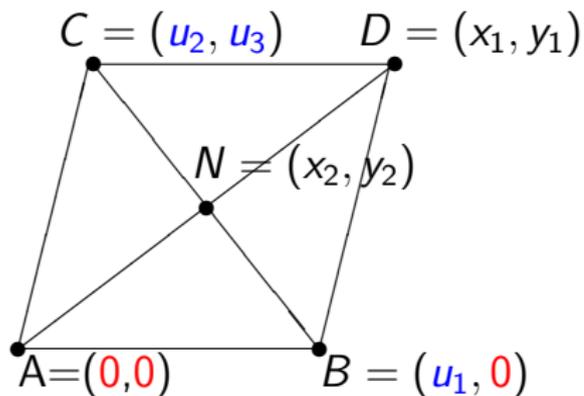


Theorem

Let $ABCD$ be a parallelogram and $N = \overline{AC} \cap \overline{BD}$.
Then N is the midpoint of \overline{AC} and \overline{BD} .

Can we automate the proof?

Application 3: Automatic Geometric Theorem Proving.



Step 1: Fix co-ordinates.

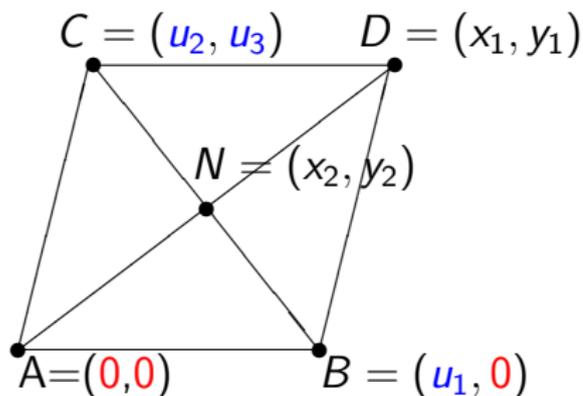
3 parameters u_1, u_2, u_3 .

4 unknowns x_1, y_1, x_2, y_2 .

Solutions are in $\mathbb{R}(u_1, u_2, u_3)$.

Step 2: Need 4 equations.

Application 3: Automatic Geometric Theorem Proving.



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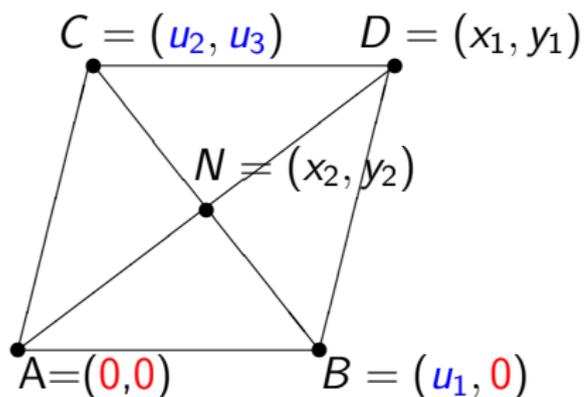
Step 2: Need 4 equations.

$ABDC$ is a parallelogram \implies the slope of $\overline{AC} = \overline{BD}$

$$\implies \frac{u_3}{u_2} = \frac{y_1}{x_1 - u_1} \implies (x_1 - u_1)u_3 = u_2y_1.$$

Similarly the slope of $\overline{AB} = \overline{CD} \implies y_1 = u_3$.

Application 3: Automatic Geometric Theorem Proving.

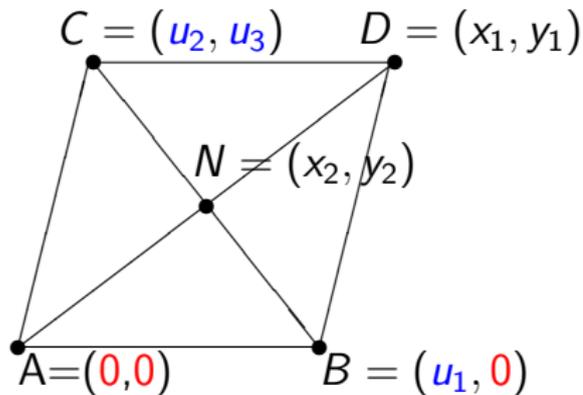


Let N be the intersection of \overline{AD} and \overline{BC} .
Hence A, N, D are co-linear \implies

$$\det \begin{pmatrix} x_2 & x_1 \\ y_2 & y_1 \end{pmatrix} = 0 \implies x_2 y_1 - y_2 x_1 = 0.$$

Similarly B, N, C are co-linear $\implies (u_1 - u_2)y_2 = u_3(u_1 - x_2)$.

Application 3: Automatic Geometric Theorem Proving.



Equations

$$h_1 = (x_1 - u_1)u_3 - u_2y_1$$

$$h_2 = y_1 - u_3$$

$$h_3 = x_2y_1 - y_2x_1$$

$$h_4 = (u_1 - u_2)y_2 - u_3(u_1 - x_2)$$

Step 2 (cont.): To prove N is the midpoint of \overline{AD} and \overline{BC}
show $\|N - A\|^2 = \|D - N\|^2$

$$\implies x_2^2 + y_2^2 = (x_1 - x_2)^2 + (y_1 - y_2)^2.$$

Similarly

$$\|N - B\|^2 = \|C - N\|^2 \implies (x_2 - u_1)^2 + y_2^2 = (u_2 - x_2)^2 + u_3^2.$$

Application 3: Automatic Geometric Theorem Proving.

$$h_1 = (x_1 - u_1)u_3 - u_2y_1$$

$$h_2 = y_1 - u_3$$

$$h_3 = x_2y_1 - y_2x_1$$

$$h_4 = (u_1 - u_2)y_2 - u_3(u_1 - x_2)$$

$$g_1 = x_2^2 + y_2^2 - (x_1 - x_2)^2 - (y_1 - y_2)^2$$

$$g_2 = (x_2 - u_1)^2 + y_2^2 - (u_2 - x_2)^2 - u_3^2$$

Step 3. Computation to prove theorem.

Let $I = \langle h_1, h_2, h_3, h_4 \rangle \in \mathbb{R}(u_1, u_2, u_3)[x_1, y_1, x_2, y_2]$.

Application 3: Automatic Geometric Theorem Proving.

$$h_1 = (x_1 - u_1)u_3 - u_2y_1$$

$$h_2 = y_1 - u_3$$

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$$g_1 = x_2^2 + y_2^2 - (x_1 - x_2)^2 - (y_1 - y_2)^2$$

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Then $g_1 \in \mathbb{V}(h_1, h_2, h_3, h_4) \iff g_1 \in \sqrt{I}$

$\iff 1 \in \langle h_1, h_2, h_3, h_4, 1 - g_1z \rangle \subset \mathbb{R}(u_1, u_2, u_3)[x_1, y_1, x_2, y_2, z]$.

Similarly verify $g_2 \in \sqrt{I}$.

What can go wrong?

Application 3: What can go wrong?

Errors: **any claim is true if** $I = \langle h_1, h_2, \dots \rangle = \langle 1 \rangle$.
 \implies check that the Gröbner basis for I is not $\{1\}$!!

Application 3: What can go wrong?

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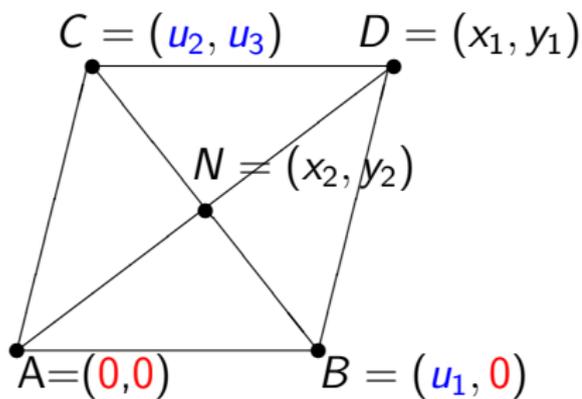
Show that working in $\mathbb{R}[u_1, u_1, u_3, x_1, y_1, x_2, y_2]$ leads to the degenerate cases $u_1 = 0$ where the theorem does not hold.

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Show that working in $\mathbb{R}[u_1, u_1, u_3, x_1, y_1, x_2, y_2]$ leads to the degenerate cases $u_1 = 0$ where the theorem does not hold.



N is the midpoint of \overline{AD}

$$\implies N = \frac{(A+D)}{2} \text{ so}$$

$$x_2 = \frac{x_1}{2}, \quad y_2 = \frac{y_1}{2} \quad !!$$

Graduate student project

- 1 Implement Buchberger's algorithm.
- 2 Study and implement the **FGLM** basis conversion.

J.C. Faugere, P. Gianni, D. Lazard, T. Mora.

Efficient computation of zero-dimensional Gröbner bases by change of ordering. *J. Symb. Comp.*, **16**, 329–344, 1993.

- 3 Show that **FGLM works** using Trinks' system.

$$\begin{aligned} &\{45p + 35s - 165b = 36, \quad 35p + 40z + 25t - 27s = 0, \\ &15w + 25ps + 30z - 18t - 165b^2 = 0, \quad -9w + 15pt + 20zs = 0, \\ &wp + 2zt = 11b^3, \quad 99w - 11sb + 3b^2 = 0\} \end{aligned}$$

Thank you for coming.

On-line course materials

www.cecm.sfu.ca/~mmonagan/teaching/MATH441