Polyomial factorization in Maple 2019.

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This is joint work with Baris Tuncer.

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Problem: Factor a in $\mathbb{Z}[x_1, x_2, \ldots, x_n]$ over $\mathbb{Q}$.

```
> a := 20*x^2*y^6*z+5*x^4*y^3*z+30*x*y^4*z^3+12*x*y^4*z^2
>     +4*x^3*y^3+3*x^3*y*z^2 +18*y^2*z^4+x^5+6*x^2*y*z^2;
```

$$20\,x^2y^6z+5\,x^4y^3z+30\,xy^4z^3+12\,xy^4z^2+4\,x^3y^3+3\,x^3yz^2+18\,y^2z^4+x^5+6\,x^2yz^2$$

```
> factor(a);
```

$$\left(5\,zxy^3+3\,yz^2+x^2\right)\left(4\,xy^3+x^3+6\,yz^2\right)$$

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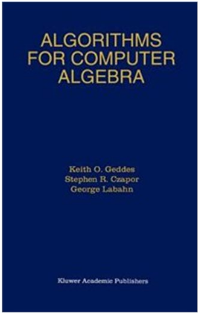
Main Tool: Multivariate Hensel Lifting (MHL)

Developed in the 1970's.
Coded in Maple by Keith Geddes in the 1980's.
Keith implemented Paul Wang's MHL which recovers the variables one at a time.
See Ch 6 of *Algorithms for Computer Algebra*, 1992.
Wang's MHL is also in Magma, Reduce, Mathematica and Singular $\Longrightarrow$ Sage.



Talk Outline

1. Wang's MHL
2. 1: Solving **multi-term** multivariate diophantine equations.
3. 2: Idea for a random polynomial time algorithm.
4. Benchmark 1
5. 3: Factors with large integer coefficients.
6. Benchmarks 2 & 3
7. Concluding Remarks

Factoring polynomials using Wang's Hensel lifting

```
> e1 := (a = f*g);
```

$$e1 := x^5 + 20\,x^2y^6z + 5\,x^4y^3z + 30\,xy^4z^3 + 12\,xy^4z^2 + 4\,x^3y^3 + 3\,x^3yz^2 + 18\,y^2z^4 + 6\,x^2yz^2 = \left(x^2 + 5\,xy^3z + 3\,yz^2\right)\left(x^3 + 4\,xy^3 + 6\,yz^2\right)$$

```
> e2 := eval(e1,z=3);
```

$$e2 := x^5 + 60\,x^2y^6 + 15\,x^4y^3 + 4\,x^3y^3 + 918\,xy^4 + 27\,x^3y + 54\,x^2y + 1458\,y^2 = \left(x^2 + 15\,xy^3 + 27\,y\right)\left(x^3 + 4\,xy^3 + 54\,y\right)$$

```
> e3 := eval(e2,y=-5);
```

$$e3 := x^5 - 1875\,x^4 - 635\,x^3 + 937230\,x^2 + 573750\,x + 36450 = \left(x^2 - 1875\,x - 135\right)\left(x^3 - 500\,x - 270\right)$$

Notes: Let h be any factor of a and let $B > \max(||h||_\infty, ||a||_\infty)$.
Multivariate Hensel Lifting (MHL) is done modulo $m = p^L > 2B$.

Wang's Multivariate Hensel Lifting (MHL) : j'th step

Input $a \in \mathbb{Z}_p[x_1, ..., x_j]$, $\alpha = (\alpha_2, \ldots \alpha_j)$, $f_{10}, \ldots, f_{r0} \in \mathbb{Z}_p[x_1, \ldots, x_{j-1}]$ s.t.
(i) $a(x_1, ..., x_{j-1}, \alpha_j) = \prod_{i=1}^{r} f_{i0}$ and
(ii) $\forall i \neq \gcd(f_{i0}(x_1, \alpha), f_{j0}(x_1, \alpha)) = 1$ in $\mathbb{Z}_p[x_1]$.

Idea: $f_i = f_{i0} + \sigma_{i1}(x_j - \alpha_j) + \sigma_{i2}(x_j - \alpha_j)^2 + \ldots$

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Initialize: $f_i \leftarrow f_{i0}$ for $1 \leq i \leq r$ and $error := a - f_1 f_2 \cdots f_r$

For $k = 1, 2, \ldots,$ while $\deg(f_1 f_2 \cdots f_r, x_j) < \deg(a, x_j)$ do

$c_k :=$ coeff($error$, $(x_j - \alpha_j)^k$)
If $c_k \neq 0$ then

Solve the MDP $\sum_{i=1}^{r} \sigma_{ik} \prod_{j \neq i} f_{j0} = c_k$ for $\sigma_{ik} \in \mathbb{Z}_p[x_1, \ldots, x_{j-1}]$.

Set $f_i \leftarrow f_i + \sigma_{ik}(x_j - \alpha_j)^k$ for $1 \leq i \leq r$.
Set $error := a - f_1 f_2 \cdots f_r$

If $error = 0$ output $(f_1, f_2, \ldots, f_r)$ else output FAIL.

Wang's MDP algorithm is $O(d^n)$ if the α_i's are non-zero.

1: Multi-term multivariate diophantine equations.

At the k'th iteration of the j'th step we must solve the multi-term MDP

$$\sum_{i=1}^{r} \sigma_{ik} \prod_{j \neq i} f_{j0} = c_k \text{ in } \mathbb{Z}_p[x_1, \ldots, x_{j-1}] \text{ s.t. } \deg(\sigma_{ik}, x_1) < \deg(\prod_{j \neq i} f_{j0}, x_1).$$

The iterative method (see [**GCL Ch. 6**]) for $r = 4$ factors

$$\begin{aligned} \sigma_1 f_2 f_3 f_4 + \sigma_2 f_1 f_3 f_4 + \sigma_3 f_1 f_2 f_4 + \sigma_4 f_1 f_2 f_3 &= c_k \\ \sigma_1 \underbrace{f_2 f_3 f_4}_{b_1} + f_1 \underbrace{\left(\sigma_2 f_3 f_4 + f_2 \underbrace{\left(\sigma_3 f_4 + \sigma_4 f_3\right)}_{\tau_2}\right)}_{\tau_1} &= c_k \end{aligned}$$

This iterative method reduces to solving $r - 1$ two term MDPs.

Idea: interpolate $\sigma_1, \sigma_2, ..., \sigma_r$ simultaneously from values. For sparse interpolation we first need to find $\mathrm{supp}(\sigma_i)$, the set of monomials $x_1^{\blacksquare} x_2^{\blacksquare} \cdots x_{j-1}^{\blacksquare}$ in σ_i .

The Taylor Coefficients : MT CASC 2016

$$f = x^3 - xyz^2 + y^3z^2 + z^4 - 2$$

Idea: $f = f_0 + \sigma_1(z-\alpha_3) + \sigma_2(z-\alpha_3)^2 + \sigma_3(z-\alpha_3)^3 + \sigma_4(z-\alpha_3)^4.$

If $\alpha_3 = 0$ then $f(z) = \underbrace{(x^3-2)}_{f_0} + \underbrace{(y^3-xy)}_{\sigma_2} z^2 + \underbrace{1}_{\sigma_4} z^4.$

If $\alpha_3 = 2$ then

$$\begin{aligned} f(z) &= \underbrace{(x^3+4\,y^3-4\,xy+14)}_{f_0} + \underbrace{(4\,y^3-4\,xy+32)}_{\sigma_1}(z-2) + \\ &\quad \underbrace{(y^3-xy+24)}_{\sigma_2}(z-2)^2 + \underbrace{8}_{\sigma_3}(z-2)^3 + \underbrace{1}_{\sigma_4}(z-2)^4 \end{aligned}$$

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$$\begin{aligned} f(z) &= \underbrace{(x^3 + 4y^3 - 4xy + 14)}_{f_0} + \underbrace{(4y^3 - 4xy + 32)}_{\sigma_1}(z-2) + \\ &\quad \underbrace{(y^3 - xy + 24)}_{\sigma_2}(z-2)^2 + \underbrace{8}_{\sigma_3}(z-2)^3 + \underbrace{1}_{\sigma_4}(z-2)^4 \end{aligned}$$

Lemma 1 [MT 2016] If α_j is chosen at random from a sufficiently large set then

$$\mathrm{Prob}[\,\mathrm{supp}(f_0) \supseteq \mathrm{supp}(f_1) \supseteq \cdots \supseteq \mathrm{supp}(f_k)\,] \text{ is high}$$

Solving $\sum_{i=1}^{r} \sigma_{ik} \prod_{j \neq i} f_{j0} = c_k$ for $\sigma_{ik} \in \mathbb{Z}_p[x_1, \ldots, x_{j-1}]$.

Let $B_i = \prod_{j \neq i} f_{j0}$ and $\sigma_{i0} = f_{i0}$.
At step k

1. Let $\sigma_{ik-1} = \sum_{j=0} a_{ij} x_1^j$ and $X_{ij} = \mathrm{supp}(a_{ij})$.

   Assume $\sigma_{ik} = \sum_{j=0} b_{ij} x_1^j$ where $b_{ij} = \sum_{k=1}^{|X_{ij}|} a_{ijk} X_{ijk}$ for unknowns a_{ijk}.

2. Pick $\beta = (\beta_2, \ldots, \beta_{j-1})$ at random from $\mathbb{Z}_p$ Needs $|X_{ij}(\beta)| = |X_{ij}|$.
3. Solve the univariate multi-term diophantine equations

$$\sum_{i=1}^{r} \sigma_{is} B_i(x_1, \beta^s) = c_k(x_1, \beta^s) \;\text{ for } 1 \leq s \leq \max |X_{ij}| \text{ for } \sigma_{is} \in \mathbb{Z}_p[x_1].$$

   Needs $\gcd(f_{i0}(x_1, \beta^s), f_{j0}(x_1, \beta^s)) = 1$ in $\mathbb{Z}_p[x_1]$ for $i \neq j$.

4. Equate coefficients and solve the shifted Vandermonde systems for a_{ijk}
   for $1 \leq i \leq r$ **for** $1 \leq j < \deg(f_i, x_1)$ solve

$$\{ b_{ij}(\beta^s) = \mathrm{coeff}(\sigma_{is}, x_1^j) \;\text{ for } 1 \leq s \leq |X_{ij}| \}.$$

Benchmark for r factors

$r/n/d/\#f_i$	Wang(MDP)	old MTSHL(MDP)	new MTSHL(MDP)
3/9/10/30	18.94(16.00)	2.26(0.60)	1.36(0.30)
4/9/15/30	OOM	104.72(23.23)	90.04(6.55)
3/9/10/50	251.20(240.77)	8.87(2.28)	4.99(0.71)
3/9/15/100	2302.7(2235.2)	122.36(**28.58**)	99.28(**8.17**)
3/11/15/100	OOM	272.78(42.74)	208.35(11.51)
3/11/10/100	515.98(424.76)	189.07(23.90)	146.80(6.25)
3/11/20/100	OOM	316.12(66.7)	256.79(19.22)

Timings in seconds for Wang's method, MTSHL from 2016 and new MTSHL.
Legend: $r = \#$factors, $n =$ number of variables, $d = \deg(f_i)$.

2: Factors with large integer coefficients

Wang's method [see **GCL** Ch. 6]

1 Factor $a(x_1, \alpha_2, \ldots, \alpha_j) = \prod_{i=1}^{r} f_i(x_1)$.
2 Pick a prime p and $L \in \mathbb{N}$ such that $p^L > 2\max(||f||_\infty, ||a||_\infty)$ where $f|a$.

Gelfond: $||f||_\infty \leq e^{d_1+d_2+\cdots+d_n}||a||_\infty$ where $d_i = \deg(a, x_i)$

3 **MHL**: Lift x_2 then x_3 etc. doing coefficient arithmetic mod p^L.
This usually means lots of multi-precision integer arithmetic.

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This usually means lots of multi-precision integer arithmetic.

We propose instead

1 same
2 same except use a machine prime p
3 **MHL**: Lift x_2 then x_3 etc. doing coefficient arithmetic mod p.

Then lift the factors mod p^2, p^3, ... to p^L stopping if error is 0.

Must solve MDPs in $\mathbb{Z}_p[x_1, \ldots, x_n]$ – in all n variables!

Consider a factor f of a with large integer coefficients.

Let p be a prime and let $f = \sum_{k=0}^{df} f_k p^k$ be a factor of $a(x_1, \ldots, x_n)$.

After MHL mod p we have computed f_0.

Example

$$\begin{aligned} f &= 2x_1 + (5 + 0 \cdot p + 2p^2)x_2 + (7 + 3p)x_3 \\ &= \underbrace{(2x_1 + 5x_2 + 7x_3)1}_{f_0} + \underbrace{3x_3}_{f_1} p + \underbrace{2x_2}_{f_2} p^2. \end{aligned}$$

$$\mathrm{supp}(f_0) \supseteq \mathrm{supp}(f_1) \not\supseteq \mathrm{supp}(f_2).$$

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$$\mathrm{supp}(f_0) \supseteq \mathrm{supp}(f_1) \not\supseteq \mathrm{supp}(f_2).$$

Theorem If p is chosen at random from $[2^{b-1}, 2^b]$ for b sufficiently large then

$$\text{Prob}[\ \mathrm{supp}(f_0) \supseteq \mathrm{supp}(f_1) \supseteq \mathrm{supp}(f_2) \supseteq \ldots \mathrm{supp}(f_{df})\ ] \text{ is high.}$$

Idea: Pick a **63 bit** prime p and assume it's true!

Algorithm p-adic lift for $r = 2$: monic case

Input : $a \in \mathbb{Z}[x_1, \ldots, x_n]$, $f_0, g_0 \in \mathbb{Z}_p[x_1, \ldots, x_n]$ such that $a = f_0 g_0$ in $\mathbb{Z}_p[x_1, \ldots, x_n]$. Also an integer lifting bound $L > 0$.

Output : $f, g \in \mathbb{Z}[x_1, \ldots, x_n]$ such that $a = fg \in \mathbb{Z}[x_1, \ldots, x_n]$ or FAIL

$(f, g) \leftarrow (f_0, g_0)$.
$error \leftarrow a - fg$.
for k from 1 to L **while** $error \neq 0$ **do**
    $c_k \leftarrow \frac{error}{p^k} \mod p$
    Solve the MDP $f_k\, g_0 + g_k\, f_0 = c_k$ for f_k and g_k in $\mathbb{Z}_p[x_1, \ldots, x_n]$
    assuming $\mathrm{supp}(f_k) \subseteq \mathrm{supp}(f_{k-1})$ and $\mathrm{supp}(g_k) \subseteq \mathrm{supp}(g_{k-1})$
    if (σ, τ)=FAIL **then return** FAIL **end if**
    $(f, g) \leftarrow (f + f_k p^k,\ g + g_k p^k)$.
    $error \leftarrow a - fg$
end for
if $error \neq 0$ **then return** FAIL **else return** (f, g) **end if**

Benchmark for p-adic lift

For $p = 2^{31} - 1$, coefficients from $(-p^l, p^l)$, L is the lifting bound.

$n/d/\#f_i$	l	L	MTSHL (MDP)	MTSHL-d (MDP) (Lift)		
5/10/300	2	5	5.866 (5.101)	0.438	(0.132)	(0.241)
5/10/500	2	5	9.265 (7.937)	1.194	(0.186)	(0.480)
5/10/1000	2	5	14.448 (12.826)	2.202	(0.264)	(1.332)
5/10/300	4	9	6.923 (6.104)	1.067	(0.156)	(0.553)
5/10/500	4	9	10.971 (9.737)	1.854	(0.219)	(1.231)
5/10/1000	4	9	16.943 (15.183)	3.552	(0.350)	(2.632)
5/10/300	4	17	8.638 (7.596)	2.553	(0.201)	(2.076)
5/10/500	4	17	13.118 (11.686)	3.101	(0.280)	(2.396)
5/10/1000	4	17	19.031 (17.225)	4.905	(0.459)	(4.032)

Timings in CPU seconds for two factors with large integer coefficients.

Compute and factor det C_n where C_n is the $n \times n$ cylic matrix. For example

$$C_4 = \begin{bmatrix} x_1 & x_2 & x_3 & x_4 \\ x_4 & x_1 & x_2 & x_3 \\ x_3 & x_4 & x_1 & x_2 \\ x_2 & x_3 & x_4 & x_1 \end{bmatrix}$$

$\det(C_4) =$
$x_1^4 - 4\,x_1^2 x_2 x_4 - 2\,x_1^2 x_3^2 + 4\,x_1 x_2^2 x_3 + 4\,x_1 x_3 x_4^2 - x_2^4 + 2\,x_2^2 x_4^2 - 4\,x_2 x_3^2 x_4 + x_3^4 - x_4^4$
$= (x_1 + x_2 + x_3 + x_4)(x_1 - x_2 + x_3 - x_4)(x_1^2 - 2\,x_1\,x_3 + x_2^2 - 2\,x_2\,x_4 + x_3^2 + x_4^2)$

n	#det	deg(f_i)	max #f_i	MTSHL	Wang	(MDP)	Magma
8	810	1,1,2,4	86	0.140s	0.096s	(52%)	0.12s
9	2704	1,2,6	1005	0.465s	0.253s	(76%)	1.02s
10	7492	1,1,4,4	715	3.03s	1.020s	(49%)	10.97s
11	32066	1,10	184756	1.33s	12.43s	(88%)	142.85s
12	86500	1,1,2,2,2,4	621	4.97s	20.51s	(65%)	7575.14s
13	400024	1,12	2704156	10.24s	212.40s	(88%)	30,871.9s
14	1366500	1,1,6,6	27132	666.0s	1364.4s	(68%)	$> 10^6$s

Table: Timings (CPU time seconds) for factoring $\det(C_n(x_n = 1))$

Concluding Remarks

Baris installed the new MTSHL code into Maple for Maple 2019.
This was done under a MITACS internship with Maplesoft.

- Michael Monagan and Baris Tuncer:
  Using Sparse Interpolation in Hensel Lifting.
  Proceedings of CASC 2016, LNCS **9890**, pp. 381–400, 2016.
- Michael Monagan and Baris Tuncer:
  Factoring multivariate polynomials with many factors and huge coefficients.
  Proceedings of CASC 2018, LNCS **11077**, pp. 319–400, 2018.
- Michael Monagan and Baris Tuncer:
  The complexity of sparse Hensel lifting and sparse polynomial factorization To appear in *J. Symbolic Computation*, 2019.

Thank you for attending!