

Lec 15B Algorithm Linear p-adic Lift

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Algorithm Linear P-adic Lifting

(10)

Inputs: p an odd prime (to recover -ve coeffs)

$f(u) \in \mathbb{Z}[x][u]$, $u_0 \in \mathbb{Z}_p[x]$ s.t. $f(u_0) = 0$ & $f'(u_0) \neq 0$

$B \in \mathbb{Z}$ s.t. $\|\tilde{u}\|_\infty < B$, a lifting bound

Output $\tilde{u} \in \mathbb{Z}[x]$ s.t. $f(\tilde{u}) = 0$ or FAIL meaning
there is no such \tilde{u} with $\tilde{u} \equiv u_0 \pmod{p}$.

(✓)

$$\tilde{u} \leftarrow u_0$$

$$d \leftarrow f'(u_0) \pmod{p} \quad (d = -zu_0)$$

for $k = 1, 2, 3, \dots$

$$e_k \leftarrow f(\tilde{u}) \quad (a - \tilde{u}^2)$$

if $e_k = 0$ output \tilde{u}

if $p^k > 2B$ output FAIL

$$t \leftarrow -\frac{e_k}{p^k} \pmod{p} \in \mathbb{Z}_p[x]$$

if $d + t \in \mathbb{Z}_p[x]$ output FAIL

$$u_k \leftarrow t/d$$

$$\tilde{u} \leftarrow \tilde{u} + u_k p^k$$

mods (Quo(t, d, x), p); in $\mathbb{Z}_p[x]$

Remark $a - \tilde{u}^2$ is expensive.

Sqrt P-adic Newton Iteration

Tell Maple to do all computations modulo p in the symmetric range.

```
> `mod` := mods;
mod := mods
```

Given the polynomial

```
> a := 49*x^4-238*x^3+513*x^2-544*x+256;
a := 49 x4 - 238 x3 + 513 x2 - 544 x + 256
```

compute $\sqrt{a(x)}$ if the sqrt exists, i.e. solve $F(u) = u^2 - a(x) = 0$ for $u(x)$. We have that $F'(u) = 2 u$. Let's first compute the sqrt modulo 5.

```
> p := 5;
p := 5
```

```
> amod5 := a mod p;
amod5 := -x4 + 2 x3 - 2 x2 + x + 1
```

We obtain that a sqrt mod 5 (by trial and error for now) is

```
> u0 := 2*x^2-2*x+1;
u0 := 2 x2 - 2 x + 1
```

```
> amod5 - ( expand(u0^2) mod p );
0
```

We need a bound on the size of the largest coefficient in the sqrt. We can use the Mignotte bound for this. Hence we must run the iteration until p^k is greater than $2 B$ where

```
> d := degree(a); B := 2^d*ceil(sqrt(d+1))*maxnorm(a);
d := 4
B := 26112
```

We are ready to go: our fo

```
> u := u0;
u := 2 x2 - 2 x + 1
```

Note that the error is calculated over Z not mod p !!

```
> e1 := a - expand(u^2);
e1 := 45 x4 - 230 x3 + 505 x2 - 540 x + 255
```

```
> e1 / 5;
9 x4 - 46 x3 + 101 x2 - 108 x + 51
```

```
> u1 := Quo(e1/5,2*u0,x,'r') mod p;
u1 := x2 + 2 x - 2
```

```
> r;
0
```

```

> u := u + u1*p;
u := 7x2 + 8x - 9

> e2 := a - expand(u^2);
e2 := -350x3 + 575x2 - 400x + 175

> e2 / 25;
-14x3 + 23x2 - 16x + 7

> u2 := Quo(e2/25, 2*u0, x, 'r') mod p;
u2 := -x + 1

> r;
0

> u := u + u2*p^2;
u := 7x2 - 17x + 16

> e3 := a - expand(u^2);
e3 := 0

```

We are done.

Consider $a(x) = 9x^2 + 18x + 24$. This polynomial obviously cannot be a perfect square because 24 is not a perfect square.

```

> a := 9*x^2+18*x+24;
a := 9x2 + 18x + 24

> amod5 := a mod 5;
amod5 := -x2 - 2x - 1

> u0 := 2*x+2;
u0 := 2x + 2

> expand( amod5 - u0^2 ) mod p;
0

> u := u0;
u := 2x + 2

> e1 := expand( a - u^2 );
e1 := 5x2 + 10x + 20

> e1/p;
x2 + 2x + 4

> d1 := Quo(e1/p, 2*u0, x, 'r') mod p;
d1 := -x - 1

> r;
-2

```

Since $r \neq 0$, \sqrt{a} is not a polynomial in $\mathbb{Z}[x]$.