

Assignment # 2 due Monday @ 11pm.  
Office hours Friday 9-11am Monday 9-11am.

## Chapter 5 Homomorphisms and Chinese Remainder Algorithm.

### 5.3 Ring Morphisms

Let  $R$  and  $S$  be two rings with identities  $1_R$  and  $1_S$ .

A function  $\phi: R \rightarrow S$  is called a ring morphism (or homomorphism) if  $\forall a, b \in R$

$$(i) \phi(a +_R b) = \phi(a) +_S \phi(b)$$

$$(ii) \quad \phi(a \cdot_R b) = \phi(a) \circ \phi(b) \quad \text{and}$$

$$(iii) \quad \phi(I_R) = I_S$$

Lemma. Let  $a \in R$ . Then

$$(iv) \quad \phi(O_s) = O_s$$

$$(v) \quad \phi(-a) = -\phi(a)$$

(vi) 'a is a unit'  $\Rightarrow \phi(a)$  is a unit

$$\text{Proof (iv)} \quad \frac{\phi(0)}{-\phi(0)} = \frac{\phi(0+0)}{-\phi(0)} \stackrel{\text{def}}{=} \frac{\phi(0)+\phi(0)}{-\phi(0)}.$$

$$(v) \quad 0 = \phi(0) = \phi(a + (-a)) \stackrel{(i)}{=} \phi(a) + \phi(-a).$$

$$\text{So } \phi(-a) = -\phi(a).$$

① The modular homomorphism  $\phi_n : \mathbb{Z} \rightarrow \mathbb{Z}_n$  where  $\phi_n(a) = a \pmod n$ .  $n > 1$

Example  $\phi_7(4 \cdot 9) = \phi_7(36) = 36 \bmod 7 = 1$ .

$$(ii) \quad \phi_7(4) \cdot \phi_7(9) = 4 \cdot 2 = 1.$$

② The evaluation homomorphism  $\phi_{x=a} : \underline{R[x]} \rightarrow R$

where  $R$  is a ring,  $a \in R$  and

$$\phi_{x=a}(f) = f(a). \quad \phi(a \cdot b) = \phi(a) \cdot \phi(b).$$

$$= -x^2 - 5x - 3$$

$$\phi_{x=a}(f) = f(a). \quad \phi(a \cdot b) = \phi(a) \cdot \phi(b).$$

Example.  $f = \frac{(x+1)(2x+3)}{3 \cdot 7} = \frac{2x^2 + 5x + 3}{8+10+3=21}$

Proof that  $\phi_n : \mathbb{Z} \rightarrow \mathbb{Z}_n$  is a ring morphism.

$$\begin{aligned} (i) \quad \phi_n(a+b) &= (\underline{a+b}) \bmod n \\ &= ((\underline{q_a \cdot n + r_a}) + (\underline{q_b \cdot n + r_b})) \bmod n \\ &= \underline{r_a + r_b} \bmod n \\ &= (\underline{a \bmod n}) + \underline{\underline{b \bmod n}} \\ &= \phi_n(\underline{a}) + \underline{\underline{\phi_n(b)}} \end{aligned}$$

$$\begin{aligned} (ii) \quad \phi_n(a \cdot b) &= \underline{a \cdot b} \bmod n \\ &= (\underline{q_a \cdot n + r_a})(\underline{q_b \cdot n + r_b}) \\ &= \underline{\underline{\phi_n(a) \cdot \phi_n(b)}} \end{aligned}$$

$$(iii) \quad \phi_n(1_{\mathbb{Z}}) = 1 \bmod \underline{n}. \quad (n > 1). \\ = 1_{\mathbb{Z}_n}$$

Theorem Let  $\phi : R \rightarrow S$  be a ring morphism.

Let  $\Theta : R[x] \rightarrow S[x]$  where  $\Theta(\sum a_i x^i) = \sum_{R[x]} \phi(a_i) x^i$   
Then  $\Theta$  is a ring morphism.  $S[x]$ .

Proof. (i). Let  $a, b \in R[x]$ .  $n = \max$  of the degrees.

$$\begin{aligned} \Theta(a+b) &= \Theta\left(\sum_{i=0}^n a_i x^i + \sum_{i=0}^n b_i x^i\right) \\ &= \Theta\left(\sum (a_i + b_i) x^i\right) \\ &= \sum \phi(a_i + b_i) x^i = \sum \underline{(\phi(a_i) + \phi(b_i))} x^i \\ &= \sum \underline{\phi(a_i)} x^i + \sum \underline{\phi(b_i)} x^i \\ &= \Theta(a) + \Theta(b). \end{aligned}$$

$$= \frac{\text{evaluate}}{\Theta(a)} + \frac{\text{evaluate}}{\Theta(b)}.$$

This extends in the obvious way to  $\Theta: R[x_1, \dots, x_n] \rightarrow S[x_1, \dots, x_n]$ .

Example.  $f = \underline{17}xy^2 + \underline{12}y^3 - 5 \in \mathbb{Z}[y][x]$

$$\phi_y(f) = 3xy^2 + 5y^3 + 2 \in \underline{\mathbb{Z}_7[y][x]}.$$

$$\begin{aligned}\phi_{y=2}(\phi_y(f)) &= 3 \cdot 2 \cdot 4 + 5 \cdot 8 + 2 \in \underline{\mathbb{Z}_7[x]} \\ &= 5x + 0.\end{aligned}$$

$$\phi_{y=2}(f) = 4 \cdot \underline{17} \cdot x + \underline{96} - 5 \stackrel{=}{=} 91 \in \mathbb{Z}[x]$$

$$\phi_y(\phi_{y=2}(f)) = 5x + 0$$

In general in a polynomial ring  $\underline{\phi_n}$  and  $\underline{\phi_{y=a}}$  commute.

Application Let  $A = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$  where  $a, b, c, d \in R$ .  $\in \mathbb{Z}[x, y]$ .

Is  $\det(A) = 0$ ?

Compute  $a \cdot d - b \cdot c$  in  $R$ . !!

Suppose  $\phi: R \rightarrow S$  as a ring morphism.

$$\begin{aligned}\phi(\det(A)) &= \phi(ad - bc) = \phi(ad) - \phi(bc) \\ &\stackrel{(ii)}{=} \phi(a) \cdot \phi(d) - \phi(b) \cdot \phi(c) = \det \left[ \begin{array}{cc} \phi(a) & \phi(b) \\ \phi(c) & \phi(d) \end{array} \right]\end{aligned}$$

If  $\det(A) = 0_R$  then  $\phi(\det(A)) = \phi(0) = 0_S$

Therefore  $\phi(\det(A)) \neq 0_S \Rightarrow \det(A) \neq 0_R$

If  $\det(A) \neq 0$  then  $\phi(\det(A))$  may be 0.

apply  $\phi$  first  
then compute  
det.

Example.  $R = \mathbb{Z}[u, v]$  Is  $\det(A) = 0$ ?

$$A = \begin{bmatrix} u & v & -uv \\ v & u & v \\ 1-u & v & u \end{bmatrix} (\phi_{u=0} \circ \phi_{v=0})(A) = \begin{bmatrix} 0 & 0 & 1 \\ 0 & 0 & 0 \\ 1 & 0 & 0 \end{bmatrix}$$

(thinking of a bigger  $A$ )

$\det = 0.$

$$(\phi_{u=1} \circ \phi_{v=1})(A) = \begin{bmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 0 & 1 & 1 \end{bmatrix}$$

$\det = -1.$

$\Rightarrow \det(A) \neq 0.$  (a proof).

$$\det(A) = (1+u-v)(u^2v+u^2-u-zv^2) \in \mathbb{Z}(u,v)$$

$\overset{u=0}{\cancel{1+u-v}}$      $\overset{u=0}{\cancel{u^2v+u^2-u}}$

If  $\det(A) \neq 0$  but  $(\phi_{u=\alpha} \circ \phi_{v=\beta})(\det(A)) = 0.$

then  $u=\alpha, v=\beta$  is a root of  $\det(A) \in \mathbb{Z}[u,v].$

What's the probability of picking a root of a polynomial?

Lemma (Schwarz-Zippel 1978)

Let  $D$  be an integral domain (e.g.  $\mathbb{Z}$ ) and  $S \subset D.$

Suppose  $f \in D[x_1, \dots, x_n]$  and  $f \neq 0.$  If  $\alpha_1, \alpha_2, \dots, \alpha_n$  are chosen at random from  $S$  then

$$\text{Prob}\left(f(\alpha_1, \alpha_2, \dots, \alpha_n) = 0 \mid \alpha \text{ is a root of } f\right) \leq \frac{\deg(f)}{|S|} = \frac{5}{10^6}$$

$S = [0, 10^6]$