

Theorem 2.4.4 (the Euclidean algorithm for polynomials)

Let $a, b \in F[x]$, F a field, $a \neq 0, b \neq 0$.

Then a and b have a unique monic gcd $g \in F[x]$.

Moreover \exists polynomials $\lambda, \mu \in F[x]$ st. $\lambda a + \mu b = g$.

Proof (existence of a gcd). Consider the sequence

$$r_0 = a, r_1 = b, r_2, \dots, r_n, r_{n+1}$$

$$(a \div b) \quad a = b q_2 + r_2 \quad \deg r_2 < \deg b$$

$$(b \div r_2) \quad b = r_2 q_3 + r_3 \quad \deg r_3 < \deg r_2$$

$$(r_{n-2} \div r_{n-1}) \quad r_{n-2} = r_{n-1} q_{n-1} + r_n \quad \deg r_n < \deg r_{n-1}$$

$$(r_{n-1} \div r_n) \quad r_{n-1} = r_n q_{n+1} + r_{n+1} \quad r_{n+1} = 0$$

Claim $r_n | a, r_n | b$, and $c|a, c|b \Rightarrow c|r_n$.

Hence $g = \text{lc}(r_n)^{-1} r_n$ is a monic gcd of $a \& b$.

(uniqueness) Let h be a gcd of $a \& b$. which is monic.

g is a common divisor $\Rightarrow gh \Rightarrow g = s \cdot h$ for $s \in F$.

h is a common divisor $\Rightarrow hg \Rightarrow h = s \cdot g$

But $\text{lc}(g) = \text{lc}(h) = 1 \Rightarrow s = 1 \Rightarrow g = h$.

(moreover) let $\lambda_0 = 1, \lambda_1 = 0, \lambda_i = \lambda_{i-2} - \lambda_{i-1} \cdot q_i \quad 2 \leq i \leq n+1$

and $\mu_0 = 0, \mu_1 = 1, \mu_i = \mu_{i-2} - \mu_{i-1} \cdot q_i \quad 2 \leq i \leq n+1$

Claim $\lambda_i a + \mu_i b = r_i$ for $0 \leq i \leq n+1$.

Proof (by induction on i)

Hence $\lambda_n a + \mu_n b = r_n$. Let $s = \text{lc}(r_n)^{-1}$

$$\Rightarrow (s\lambda_n) a + (s\mu_n) b = s r_n = g$$

$\underset{x}{\lambda} \quad \underset{\mu}{\mu}$