

Assignment #4 due tonight @ 11pm.

Assignment #5 due in two weeks.

CLO 1.2 Affine Varieties in k^n where k is a field.

Let $f_1, f_2, \dots, f_s \in k[x_1, \dots, x_n]$, $s > 0$.

Let $V(f_1, \dots, f_s) = \{(a_1, \dots, a_n) \in k^n : f_i(a_1, \dots, a_n) = 0 \ \forall 1 \leq i \leq s\}$.

This is the set of solutions of the polynomial system $\{f_1 = 0, \dots, f_s = 0\}$.

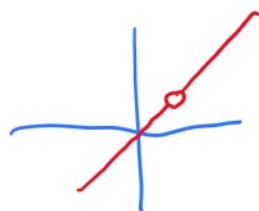
E.g. $V(0) = k^n$, $V(1) = \{\emptyset\}$.

Def. Let $S \subseteq k^n$. S is an affine variety $\Leftrightarrow \exists f_1, \dots, f_s \in k[x_1, \dots, x_n]$
s.t. $S = V(f_1, \dots, f_s)$.

Examples in \mathbb{R}^2

$$\text{A circle } \text{--- } x^2 + y^2 = 1$$

$$= V(x^2 + y^2 - 1)$$



is not an affine variety.

Examples in \mathbb{R}^3

$$V(x, y) = \{(0, 0, t) : t \in \mathbb{R}\} = z \text{ axis.}$$

$$V(xz, yz) \Rightarrow xz = 0 \text{ and } yz = 0$$

$$\begin{array}{c} \text{z axis} \Rightarrow z = 0 \text{ or } x = 0 \text{ and } y = 0. \\ \text{xy plane} = V(z) \cup V(x, y) \\ \text{one irreducible varieties.} \end{array}$$

Can such a decomposition $V(xz, yz) = V(z) \cup V(x, y)$ be computed algorithmically? Yes. GB + factor over k .
Polynomial time? In $s, \deg(f_i), n$? No.

CLO 1.4 Ideals in $k[x_1, \dots, x_n]$

Def. (Ideal) Let R be a comm. ring. E.g. \mathbb{Z} , $k[x_1, \dots, x_n]$.

A subset I of R is an ideal of R if

- (i) $0 \in I$
- (ii) $f, g \in I \Rightarrow f+g \in I$
- (iii) $f \in I, h \in R \Rightarrow fh \in I$.

What's the smallest ideal in R ? $\{0\}$ Trivial Ideals.
What's the largest ideal in R ? R .

Example. $R = \mathbb{Z}$. $I = 2\mathbb{Z} = \{0, \pm 2, \pm 4, \dots\}$ is an ideal.

Ideals in $k[x_1, \dots, x_n] = R$

Let $f_1, \dots, f_s \in R$. Define

$$\langle f_1, \dots, f_s \rangle = \{ h_1 f_1 + h_2 f_2 + \dots + h_s f_s : h_i \in R \}.$$

Lemma. $I = \langle f_1, \dots, f_s \rangle$ is an ideal of R .

Proof. (i) $h_i = 0$: $0 \cdot f_1 + 0 \cdot f_2 + \dots + 0 \cdot f_s = 0 \Rightarrow 0 \in I$.

(ii). Let $f, g \in I \Rightarrow f = \sum_{i=1}^s h_i f_i$ and $g = \sum_{i=1}^s t_i f_i$
where $h_i, t_i \in R$.

$$f+g = \sum h_i f_i + \sum t_i f_i = \sum (h_i + t_i) f_i \stackrel{R}{\Rightarrow} f+g \in I.$$

(iii) Let $f \in I, h \in R \Rightarrow f = \sum h_i f_i$
 $\Rightarrow f \cdot h = (\sum h_i f_i) \cdot h = \sum (h_i \cdot h) f_i \stackrel{R}{\in} I$.

Exercise. If I and J are ideals of R .

Prove or disprove $I \cap J$ and $I \cup J$ is an ideal of R .

Let $I = \langle f_1, \dots, f_s \rangle$ be an ideal of $k[x_1, \dots, x_n]$.

We say $\{f_1, \dots, f_s\}$ generates I or is a basis for I .

Bases are not unique.

Ex. Let $I = \langle x, y \rangle$ in $k[x, y]$.

Then $2x+y \in I$ because $h_1=1, h_2=1$, $h_1 \cdot x + h_2 \cdot y \in I$.

Consider $J = \langle x, y, 2x+y \rangle$. I claim $I=J$. Proof.

($I \subset J$): $f \in I \Rightarrow f = A \cdot x + B \cdot y = \underline{A} \cdot x + \underline{B} \cdot y + \underline{0} \cdot (2x+y) \in J$.

($J \subset I$): $f \in J \Rightarrow f = A \cdot x + B \cdot y + C(2x+y)$

$$= (A+C)x + (B+C)y \in I.$$

Lemma. Let $I = \langle f_1, \dots, f_s \rangle$ and $J = \langle g_1, \dots, g_t \rangle$.

If $f_i \in J$ and $g_i \in I$ then $I=J$.

($I \subset J$) Let $f \in I \Rightarrow f = \sum h_i f_i \Rightarrow f = \sum_{i=1}^s h_i \left[\sum_{j=1}^t a_{ij} g_j \right]$

$$\text{But } f_i \in J \Rightarrow f_i = \sum_{j=1}^t a_{ij} \cdot g_j = \sum_{j=1}^t \left[\sum_{i=1}^s a_{ij} h_i \right] g_j \in J.$$

($J \subset I$). Same argument.

Ex. $I = \langle \frac{f_1}{x+y}, \frac{f_2}{x+z}, \frac{f_3}{y-z} \rangle$.

Notice $f_3 = f_1 - f_2 \Rightarrow f_3 \in \langle x+y, x+z \rangle = J$.

Thus $I = \langle x+y, x+z \rangle$.

$$h_1=1 \quad h_2=-1 \\ f_3 = f_1 - f_2 \in J.$$

Def. We say $B = \{f_1, \dots, f_s\}$ is a minimal basis for $I = \langle f_1, \dots, f_s \rangle$ if $B \setminus \{f_i\}$ is not a basis for I .

Claim $\{x+y, y-z\}$ is a minimal basis for $\langle x+y, y-z \rangle$.

Proof. Is $I = \langle x+y \rangle$? No $y-z \notin \langle x+y \rangle = \{h_1 \cdot (x+y) : h \in R\}$, as $y-z$ is not a multiple of $x+y$.

Is $I = \langle y-z \rangle$? No $x+y \notin \langle y-z \rangle$.