

The division algorithm

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CLO 2.3 A division algorithm for $k[x_1, \dots, x_n]$.

Given $f_1, \dots, f_s \in k[x_1, \dots, x_n] \setminus \{0\}$, $f \in k[x_1, \dots, x_n]$, to divide $f \div \{f_1, \dots, f_s\}$ we want to write

$$f = a_1 f_1 + \dots + a_s f_s + r \quad \text{for } a_1, \dots, a_s, r \in k[x_1, \dots, x_n].$$

\uparrow quotients \uparrow remainder.

Context $I = \langle f_1, \dots, f_s \rangle$. Is $f \in \langle f_1, \dots, f_s \rangle$.
If $r=0 \Rightarrow f \in I$.

Example Suppose $f_1 = xy+1$, $f_2 = 1+y$, $f = -x+xy^2$.
Suppose we use lex with $x > y$. $\Rightarrow f = xy^2 - x$

$$\begin{aligned} f_1 &= xy+1 \\ f_2 &= y+1 \end{aligned}$$

$$\begin{aligned} a_1 &= y \\ a_2 &= -1 \end{aligned} \quad f = -x+1$$

$$\begin{array}{r} \overline{xy^2 - x} = f = p_1 \\ -(xy^2 + y) \end{array}$$

$$-x-y = p_2$$

$$\downarrow \quad -y = p_3$$

$$-1 \cdot f_2 - (-y-1)$$

$$\downarrow \quad 1 = p_4$$

$f = a_1 f_1 + a_2 f_2 + r$?
No term in r is divisible by $\text{LT}(f_i)$.

Notice.

$$xy^2 > x > y > 1$$

Does $r \neq 0$ mean

$f \notin \langle f_1, f_2 \rangle$?

$$\begin{aligned} a_2 &= xy-x \\ a_1 &= 0 \end{aligned} \quad r=0.$$

$$\begin{aligned} f_2 &= y+1 \\ f_1 &= xy+1 \end{aligned}$$

$$\begin{array}{r} \overline{xy^2 - x} = f \\ -(xy^2 + xy) \\ -xy-x \\ -xf_1 - (-xy-x) \end{array}$$

$f \in \langle f_1, f_2 \rangle$.

The output depends on the order of f_i 's.
 If $f \in \langle f_1, \dots, f_s \rangle$ division may not produce a 0 remainder.
 The problem is the basis $\{f_1, \dots, f_s\}$ for I not $R \models \text{alg.}$

Proof of termination

Claim Each time round the loop $\underline{\text{LM}(P_{\text{new}})} < \underline{\text{LM}(P_{\text{old}})}$ or $P_{\text{new}} = 0$.

CASE 1. Is $p - \text{LT}(p) = 0$ or $\underline{\text{LM}(P - \text{LT}(p))} < \underline{\text{LM}(P)}$. ~~($x^2 + xy + y^2$)~~

CASE 2. Is $\underline{\text{LM}(P - t f_i)} < \underline{\text{LM}(P)}$ or $\underline{P - t f_i} = 0$.

$$\begin{aligned} & \underline{\text{LM}(P - t f_i)} && t \\ &= \underline{\text{LM}(\text{LT}(p) + (P - \text{LT}(p)))} - \frac{\underline{\text{LT}(p)}}{\underline{\text{LT}(f_i)}} \cdot (\underline{\text{LT}(f_i)} + (f_i - \underline{\text{LT}(f_i)})) && > \\ &= \underline{\text{LM}(\text{LT}(p) + P - \text{LT}(p))} - \left[\underline{\text{LT}(p)} - t(f_i - \underline{\text{LT}(f_i)}) \right]. && \quad \text{LT}(p) \text{ by (i) } \\ & \quad P - \quad && < \underline{\text{LT}(p)} \end{aligned}$$

Letting P_1, P_2, \dots denote the values of P at the i^{th} step of $R \models \text{alg.}$, i.e., $P_1 = p$. then

$$\text{LM}(p_1) > \text{LM}(p_2) > \text{LM}(p_3) > \dots.$$

Since $>$ is a well ordering, Lemma 2 says such a strictly decreasing sequence cannot continue indefinitely, hence $P=0$, and $R \models \text{alg.}$ terminates.